

## Internet Sink Deployments: On the Design and Use of Internet Sinks for Network Abuse Monitoring

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# iSink Overview

- „ Definitions, Approach
- „ Architecture
- „ Performance
- „ Real World Observations
- „ Performance Evaluation
- „ Active Responses, Tarpit Effectiveness
- „ Results, Discussion
- „ Future Work

# iSink Definitions

## • Internet Sink:

- A system, either passive or active, to which IP traffic is diverted
- Includes blackhole/sinkhole routers, tarpits

## • Network Abuse:

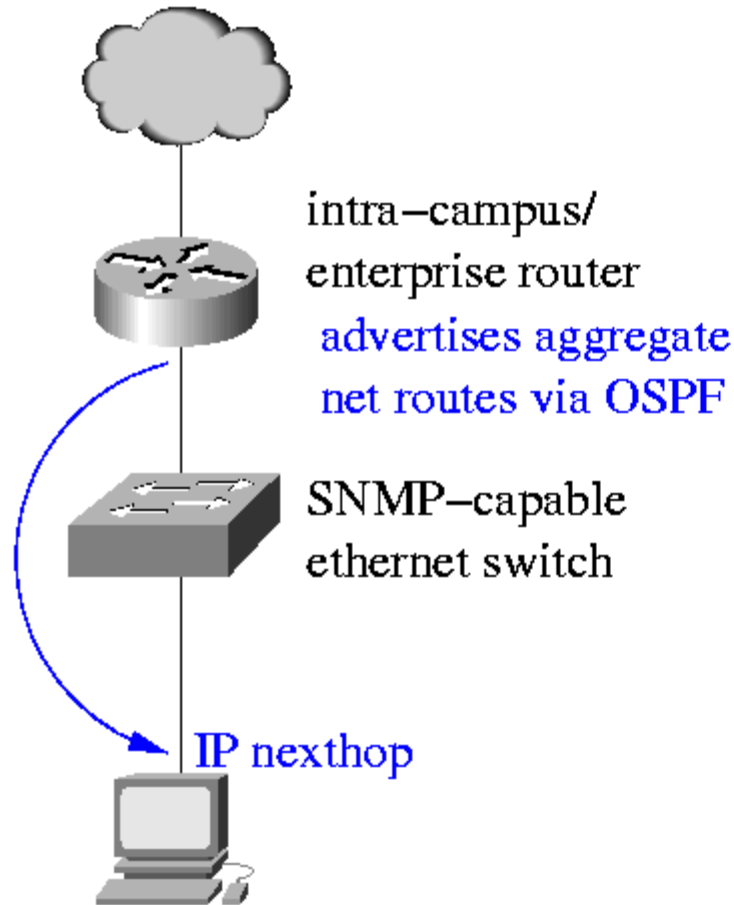
- Notion is broadly defined, ultimately determined by local policies
- We focus on intrusion attempts and attack activity

# iSink Approach

- Monitor activity of "unused" IP addresses or transport endpoints:
  - This traffic is not often monitored
  - Limited false positives:
    - misconfigurations
    - typos
  - System may be configured to respond without interfering with existing services, ie. an active sink
  - Lower traffic rates
    - sampling not necessarily required
  - Anomalies are unobscured by production traffic

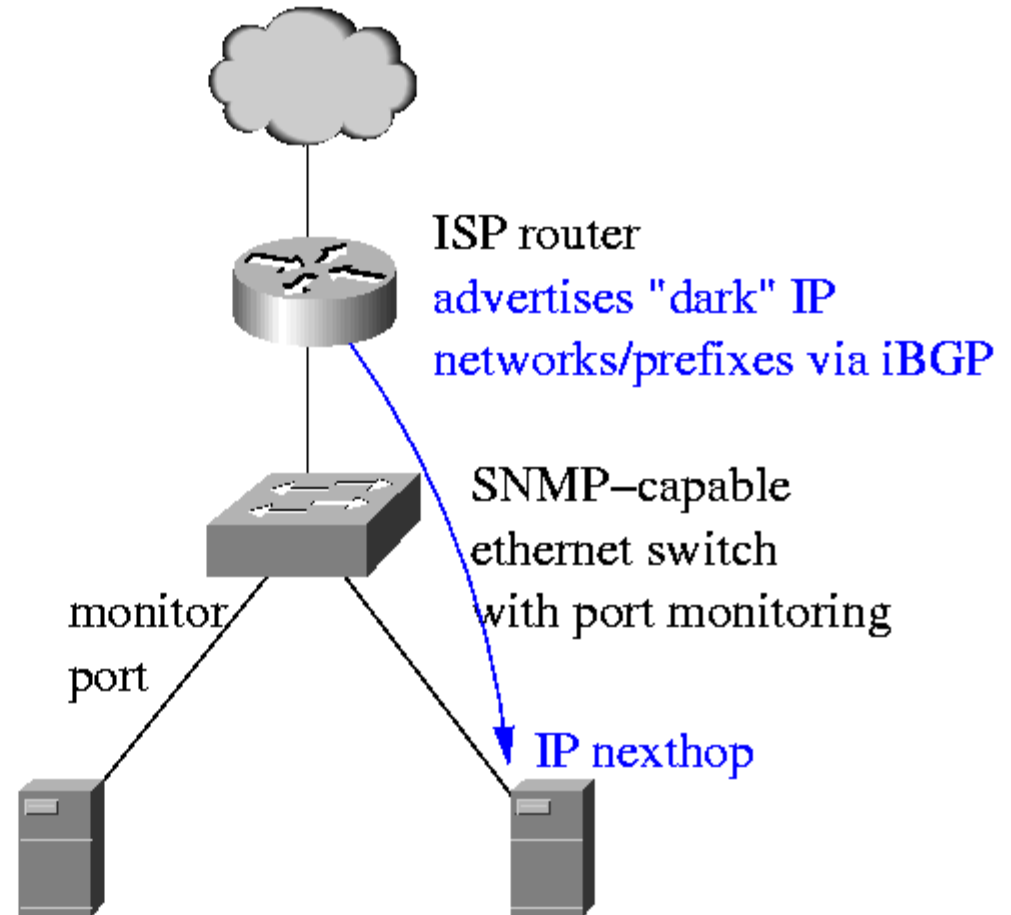
# iSink Architecture

## Enterprise Backbone/Core



argus and LaBrea  
passive/active sink

## ISP Backbone/Core

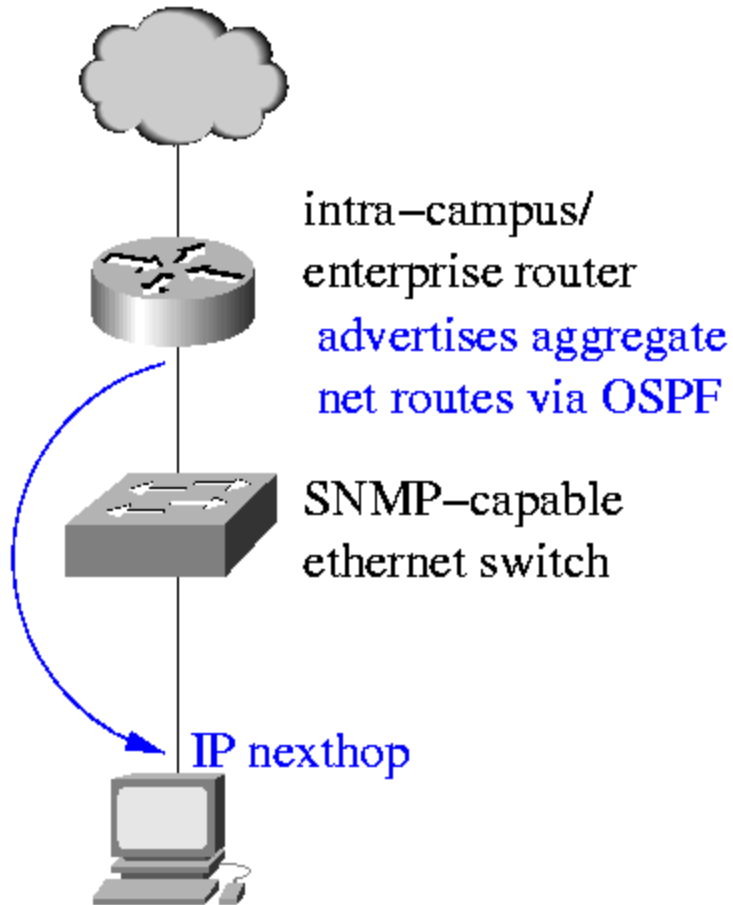


argus-based  
passive sink

click-based  
active sink

# Enterprise iSink Architecture

Enterprise Backbone/Core

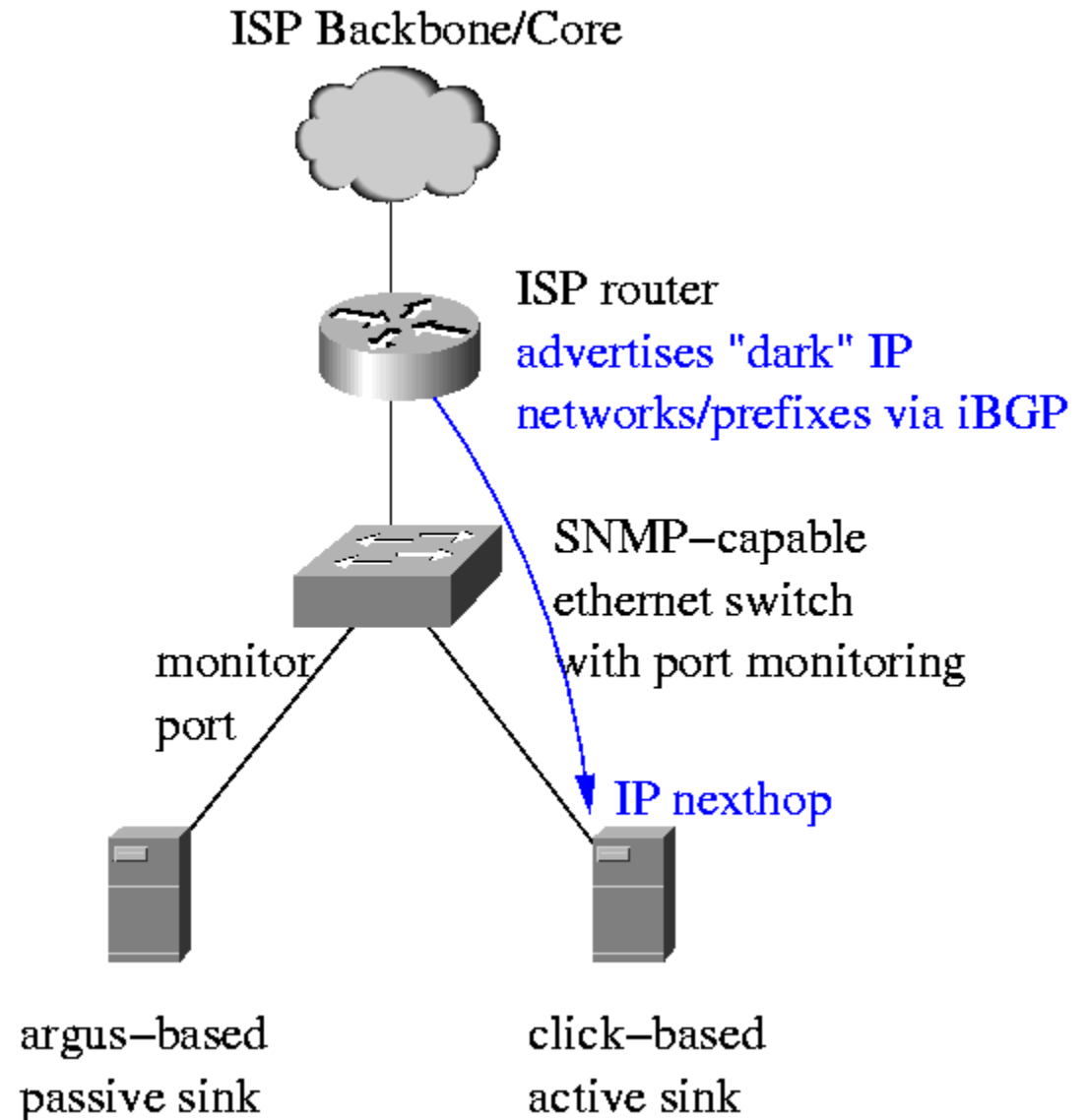


argus and LaBrea  
passive/active sink

- Receives traffic destined for ~100,000 IP addresses across four class B networks
  - In /25 - /22 blocks
  - These are the "holes" between campus subnets
- Actively responds using the LaBrea tarpit software

# ISP iSink Architecture

- „ Receives unsolicited traffic for 16 million IP addresses, one entire class A network
- „ Actively responds in some subnets:
  - SYNACKer
  - Linux
  - Windows
  - Solaris

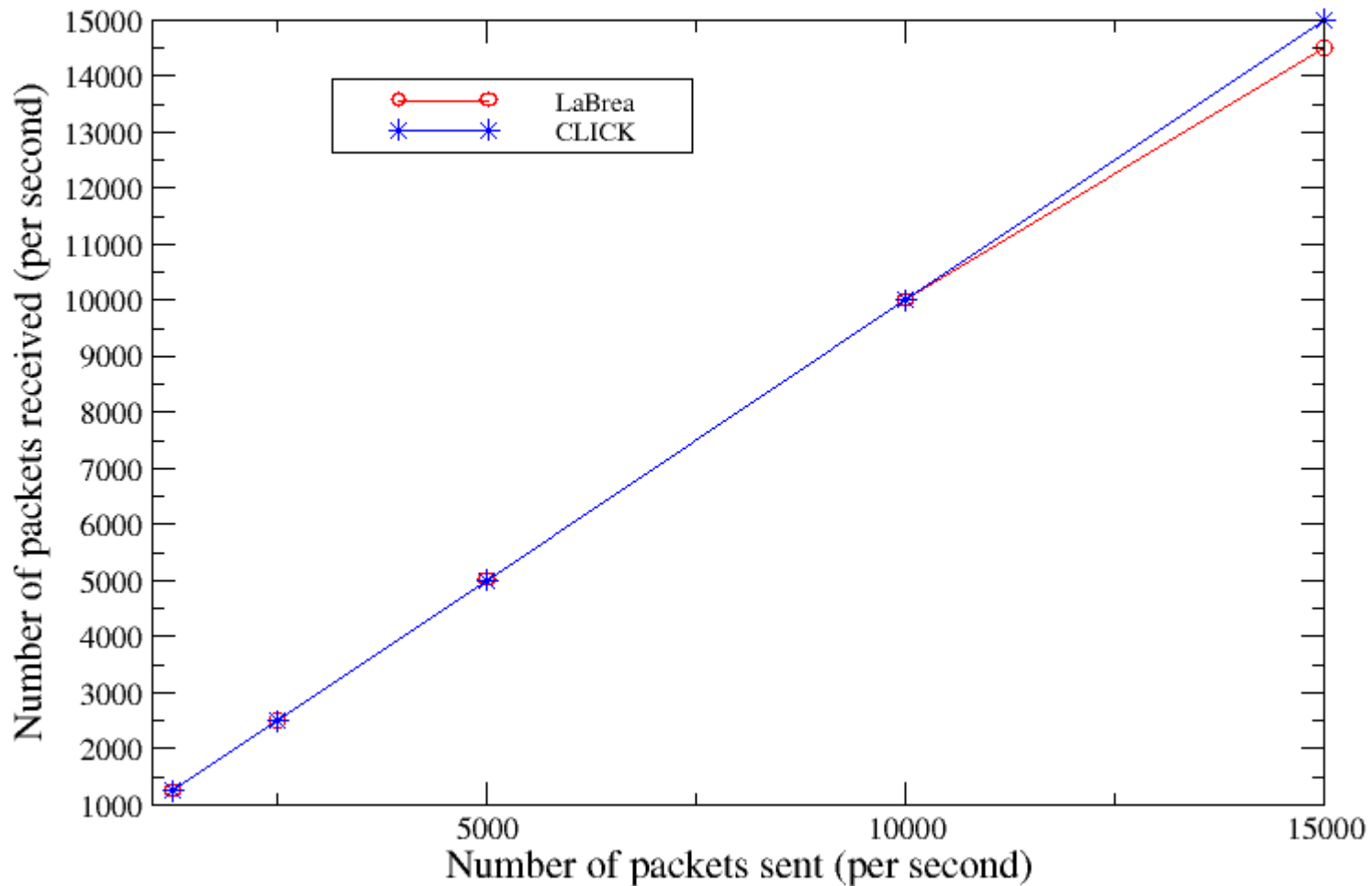


# iSink Performance

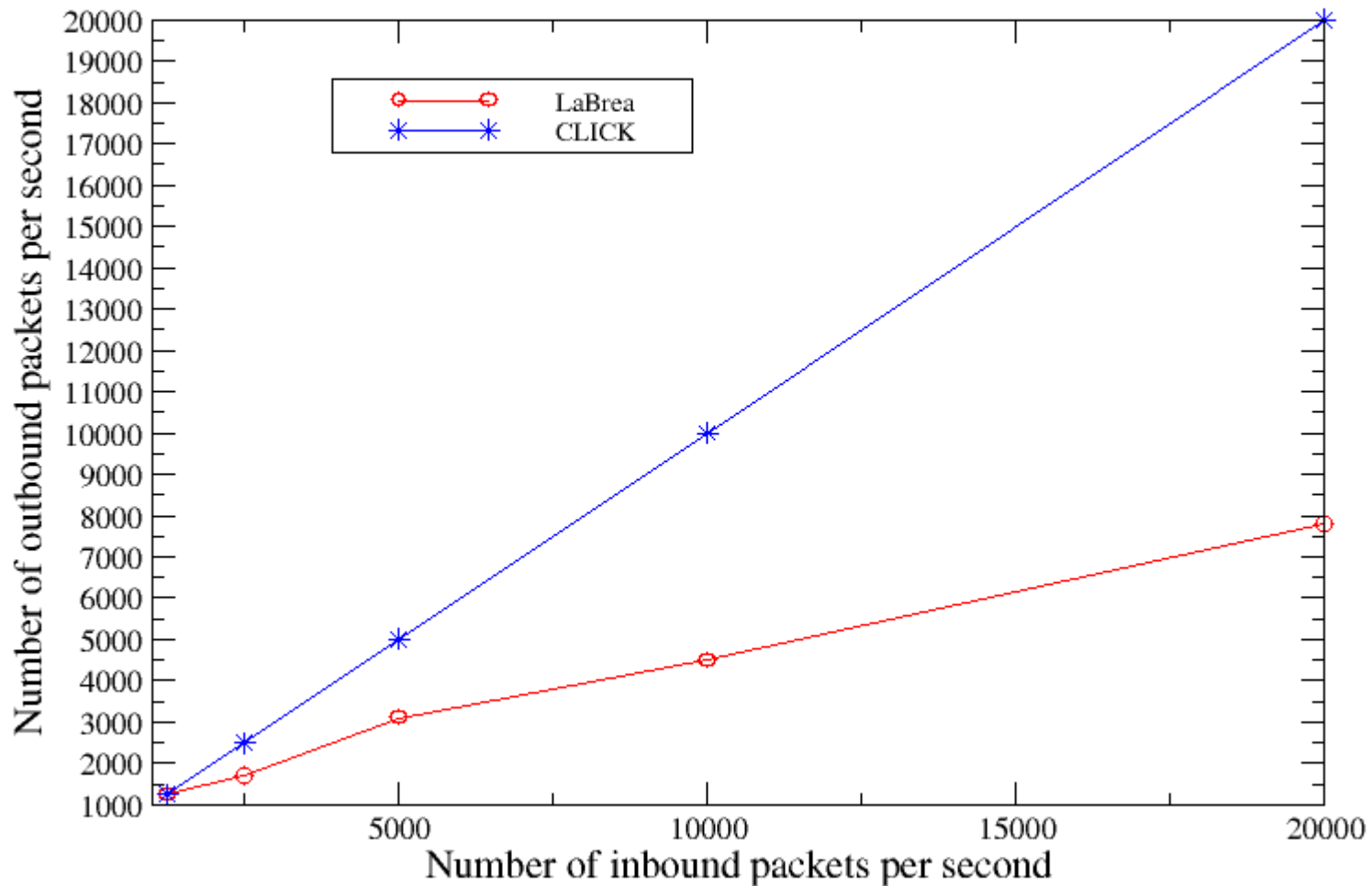
- We Considered Extensibility and Scalability
- Compare LaBrea and click-based implementations on Pentium 4 machines:
  - determine TCP and UDP response capacity
  - ARP
- Click handled 20k TCP SYNs per second
- LaBrea handled 2k TCP SYNs per second
- Both can handle ~15k UDP packets per second



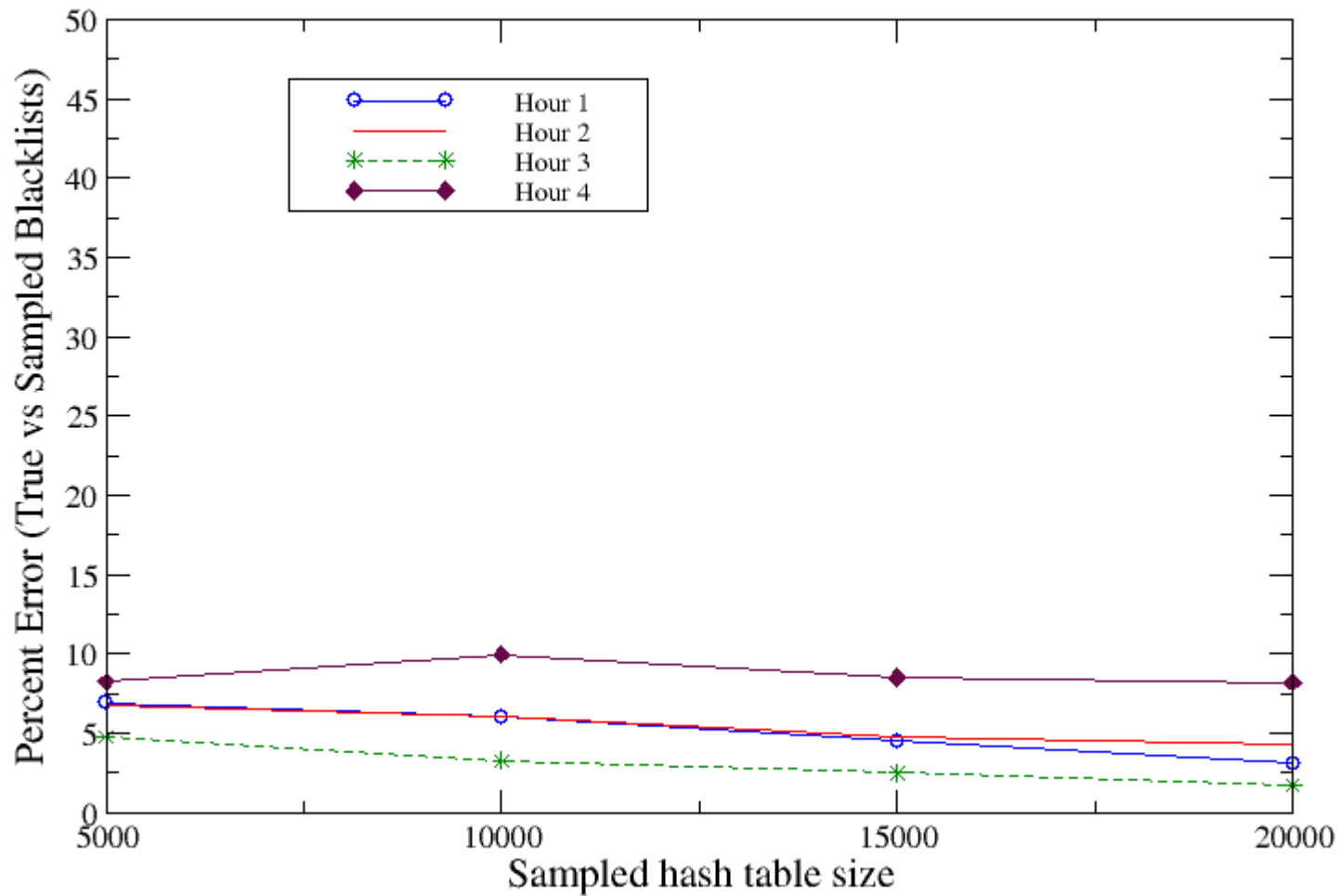
# iSink UDP Performance



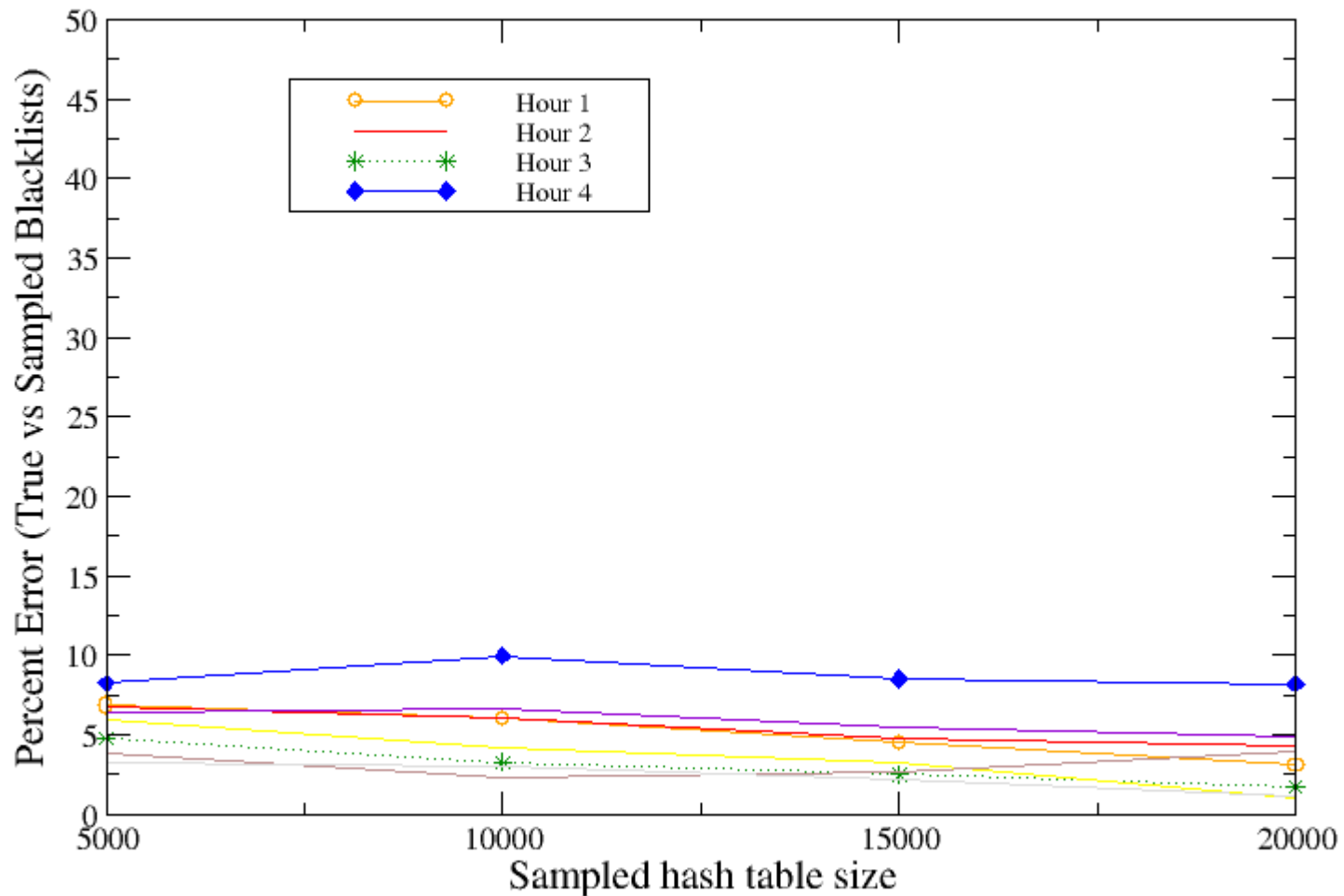
# iSink TCP Performance



# iSink Sampling 1/300



# iSink Sampling 1/100



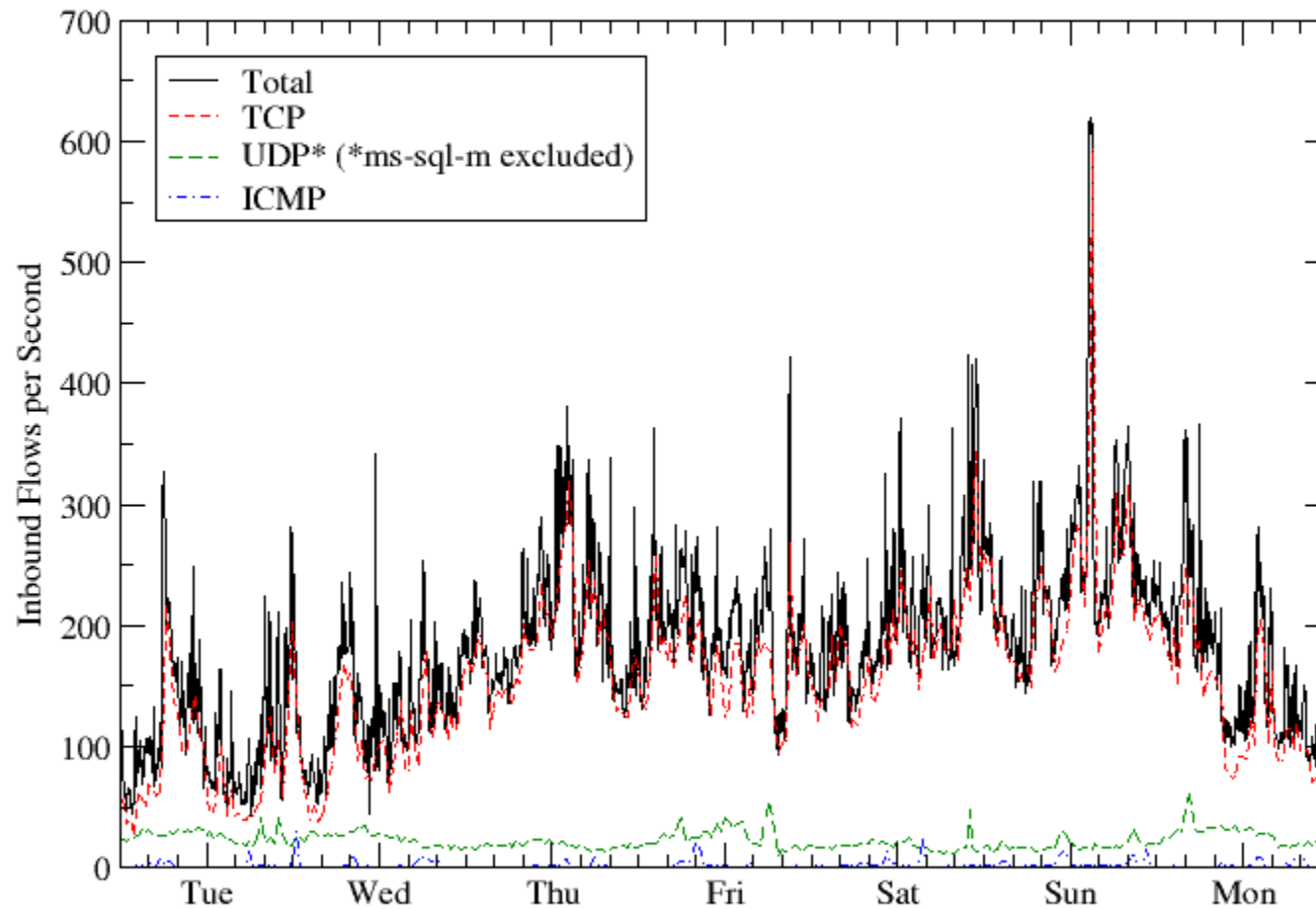
# iSink Observations

- Campus iSink received:
  - Traffic from local Network Management Systems such as ping and SNMP query attempts
  - Traffic from misconfigured campus hosts, such as DNS queries to addresses which were presumably DNS servers at some time
    - the IP address space's history matters
  - Traffic from malicious probes or worms with an affinity for the local network or presumed subnet

# iSink Observations

## Inbound Flows by IP Protocol

WAIL Campus Network Sink (~100K Addresses)



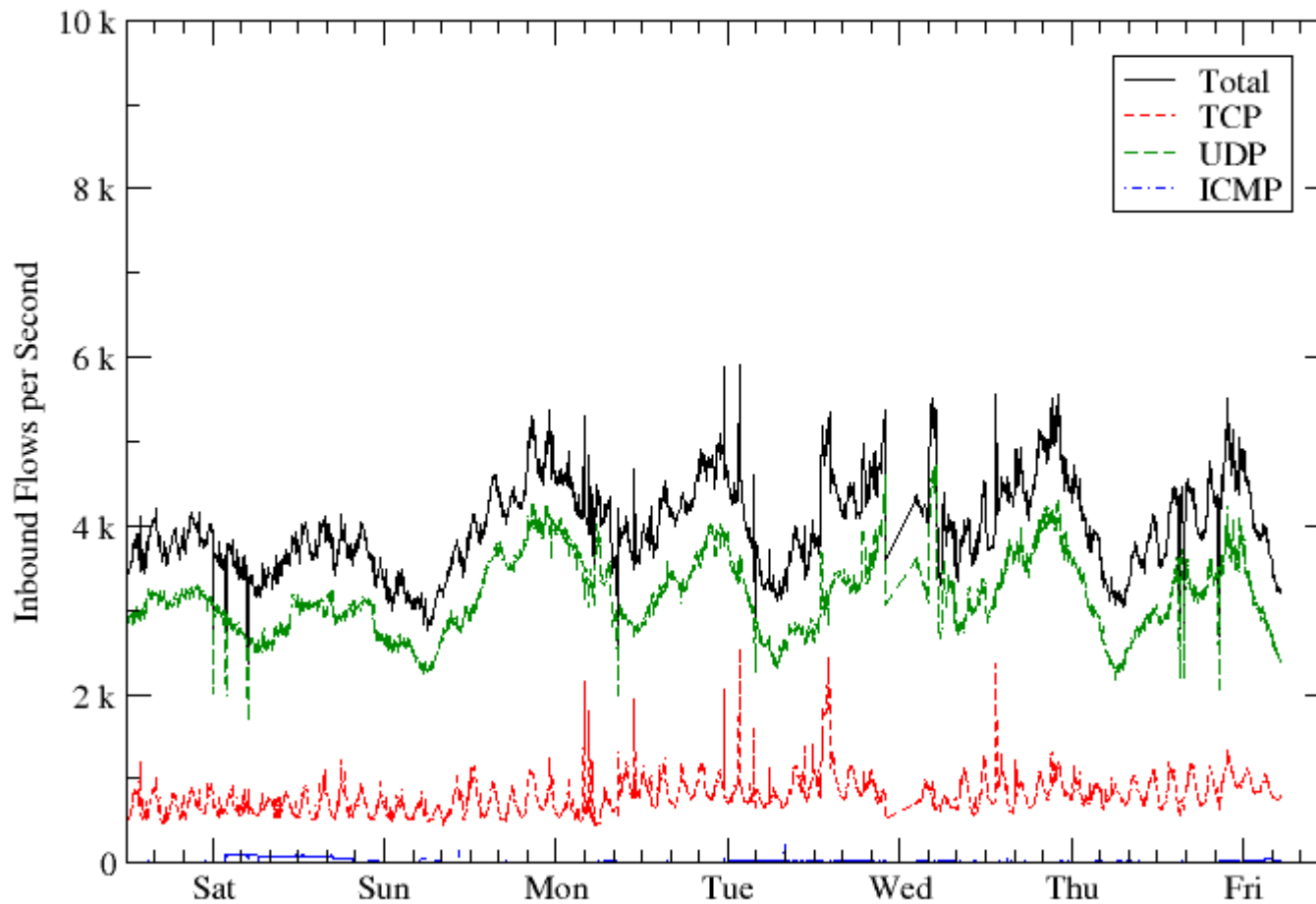
# iSink Observations

- Service-Provider iSink observed:
  - Average inbound rates of 5k packets per second and over 5Mbits per second
  - Not all networks respected more-specific routes in the global BGP, local policies sometimes prevent traffic from routing to iSinks
  - Unsolicited traffic dominated by UDP, due to popular probed services at the time
  - Significant amount of backscatter: traffic from Internet hosts presumed to be under attack from malicious parties forging the source IP address

# iSink Observations

## Inbound Flows by IP Protocol

WAIL Class-A Network Sink (16M Addresses)





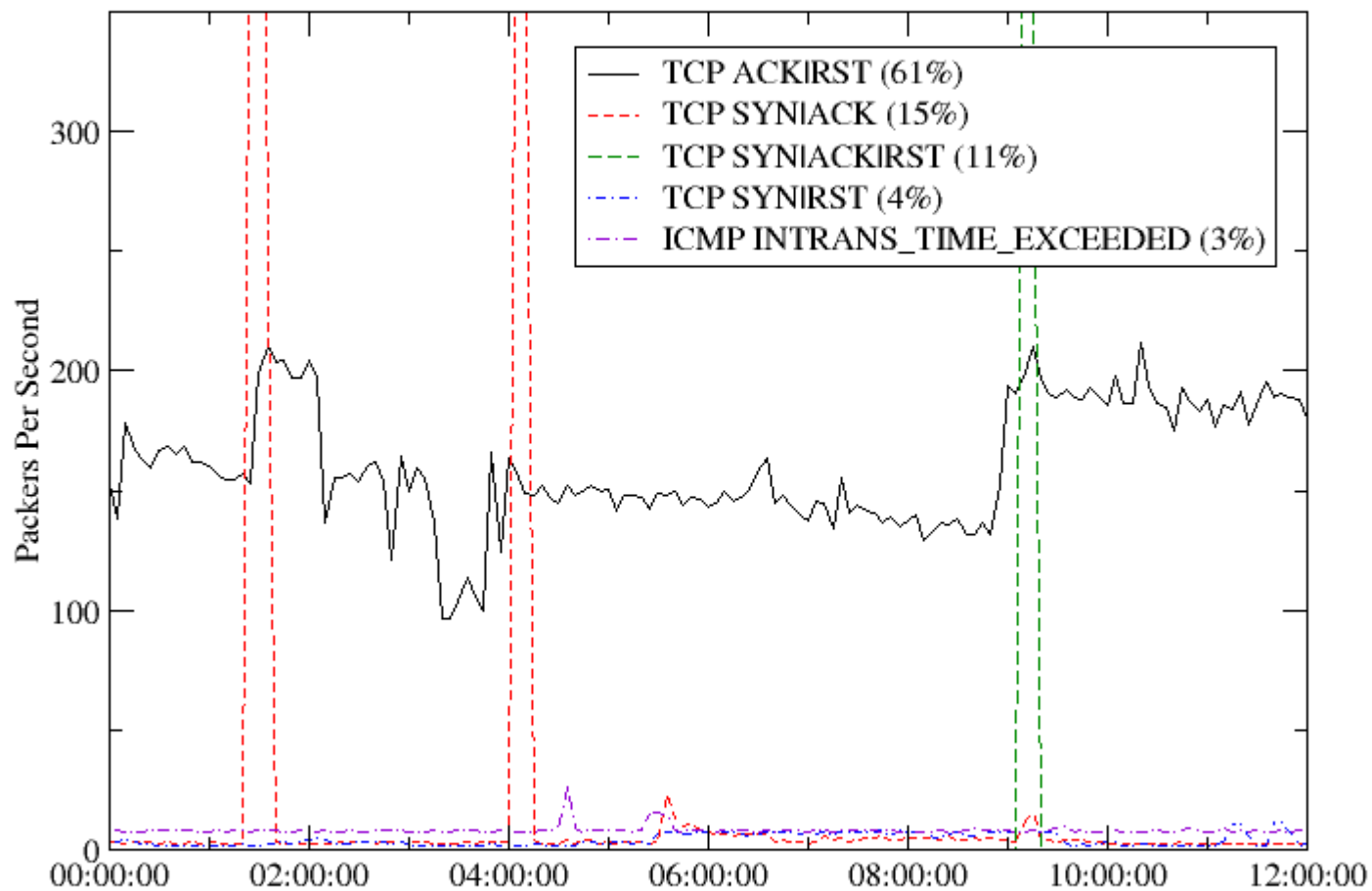
# iSink Top Services

<u>Service</u>	<u>Inbound flows per second</u>
UDP netbios- ns destination	1932
UDP ms- sql- m destination	1187
http destination	197
netbios- ssn destination	133
microsoft- ds destination	115
smtp destination	67
http source	44
http destination	11
ms- sql- s destination	10
telnet destination	2

# iSink Observations

## Inbound Backscatter Packets

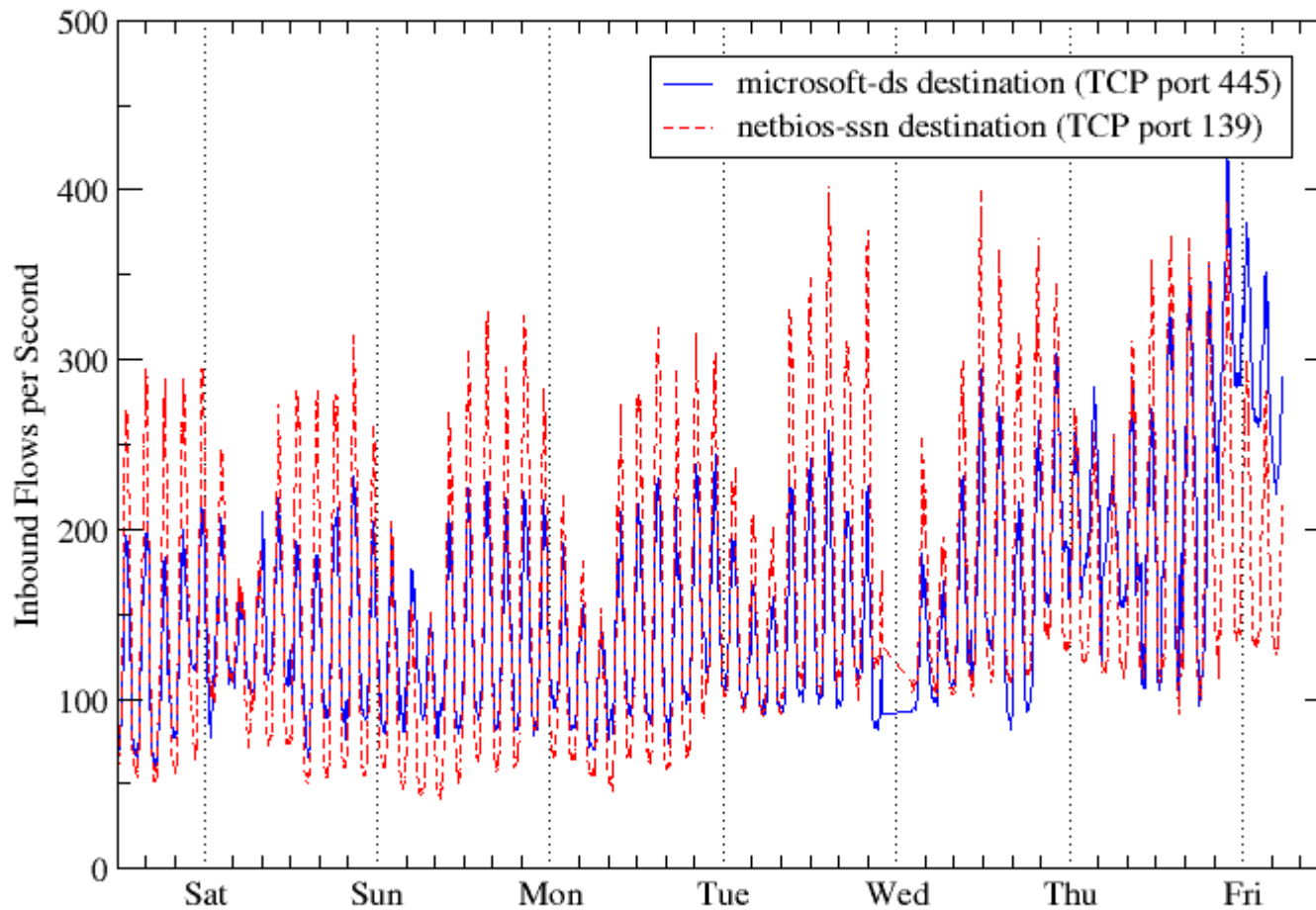
WAIL Class-A Network Sink (16M Addresses), 12 hours



# iSink Observations

Periodic Service Probing, period =  $\sim 2.67$  hours

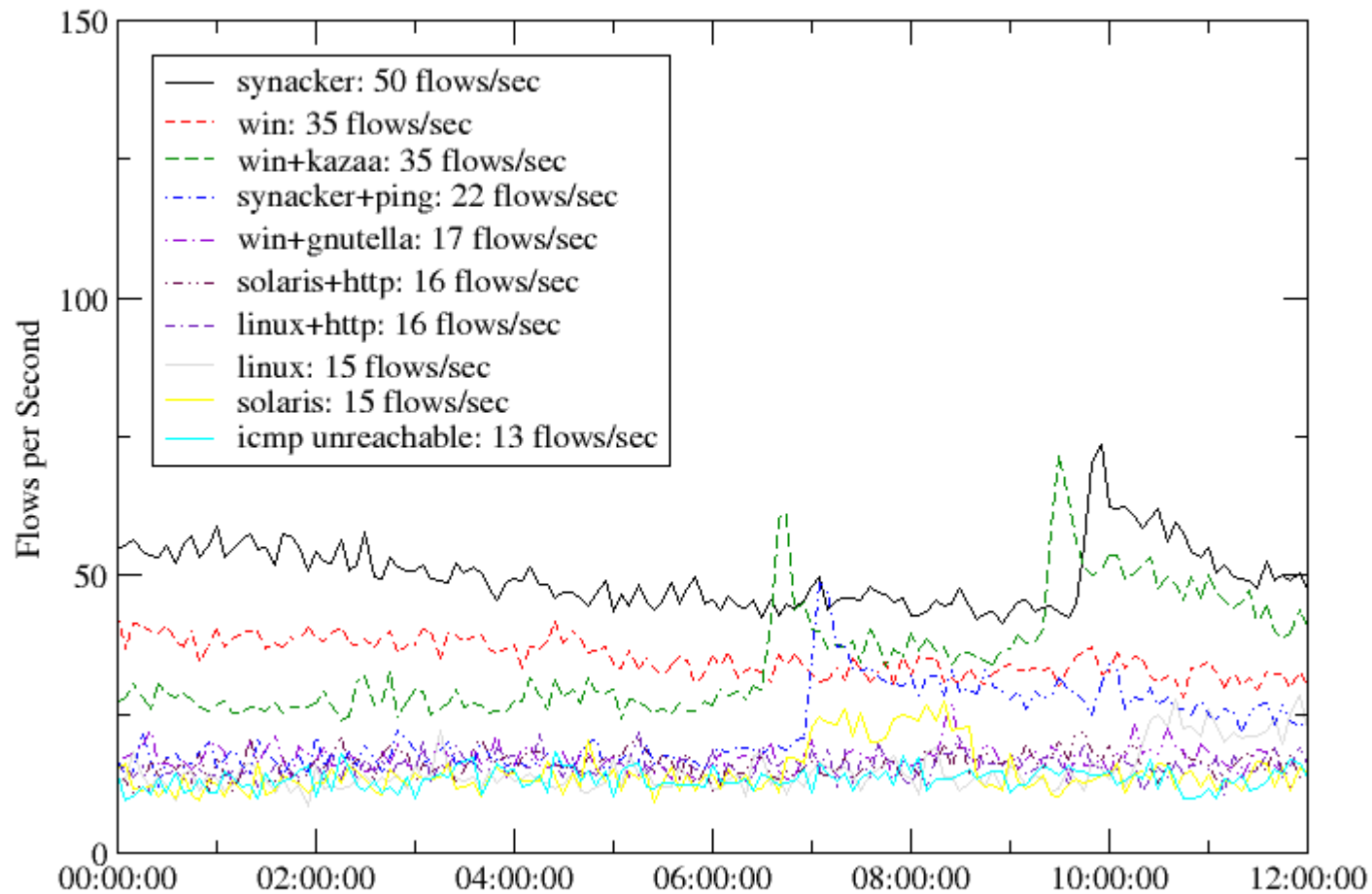
WAIL Class-A Network Sink (16M Addresses)



# iSink Active Responses

## Inbound Flows, Differentiated Responses

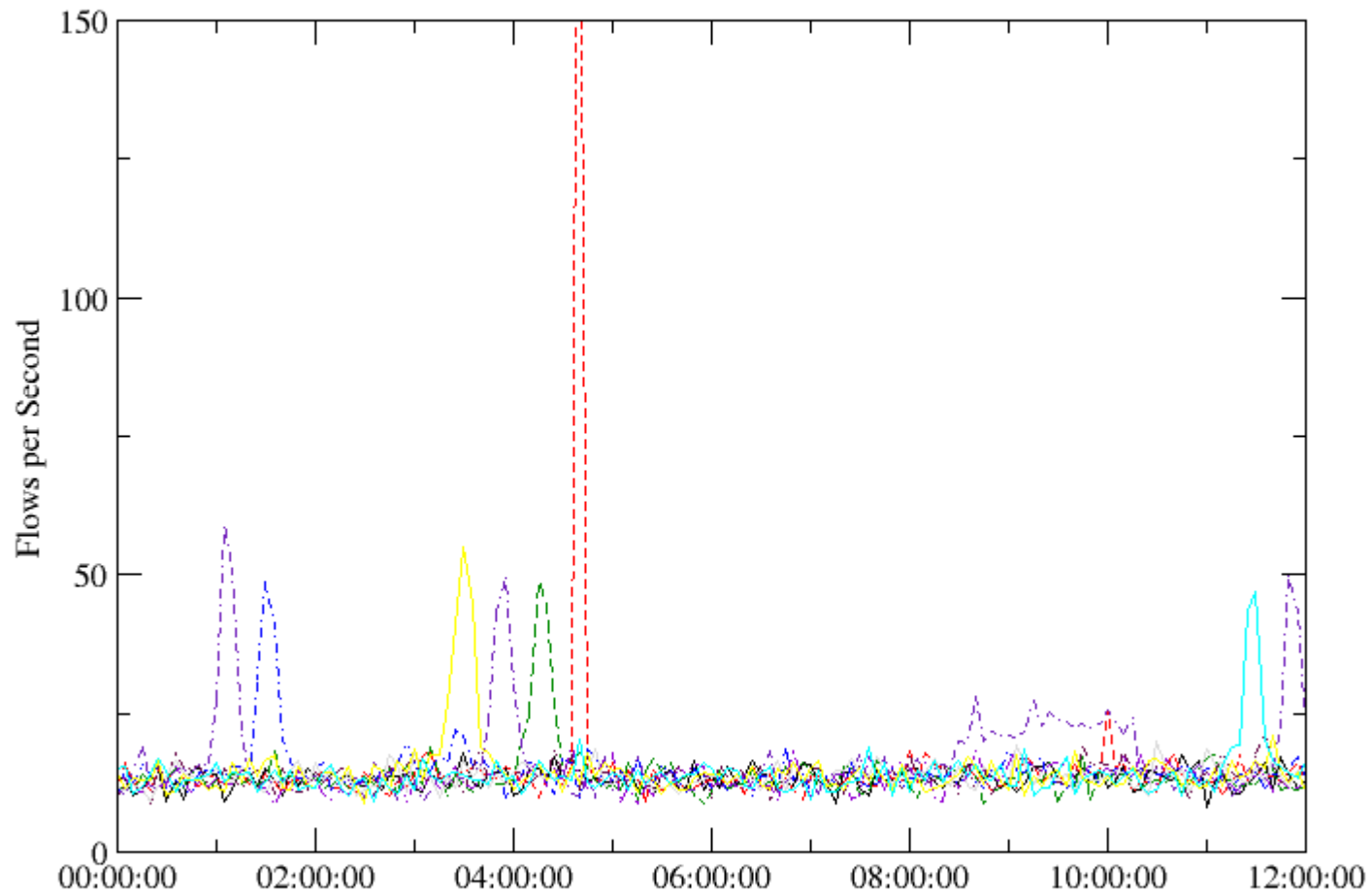
WAIL Class-A Network Sink Active Responses on 10 "/16" Subnets



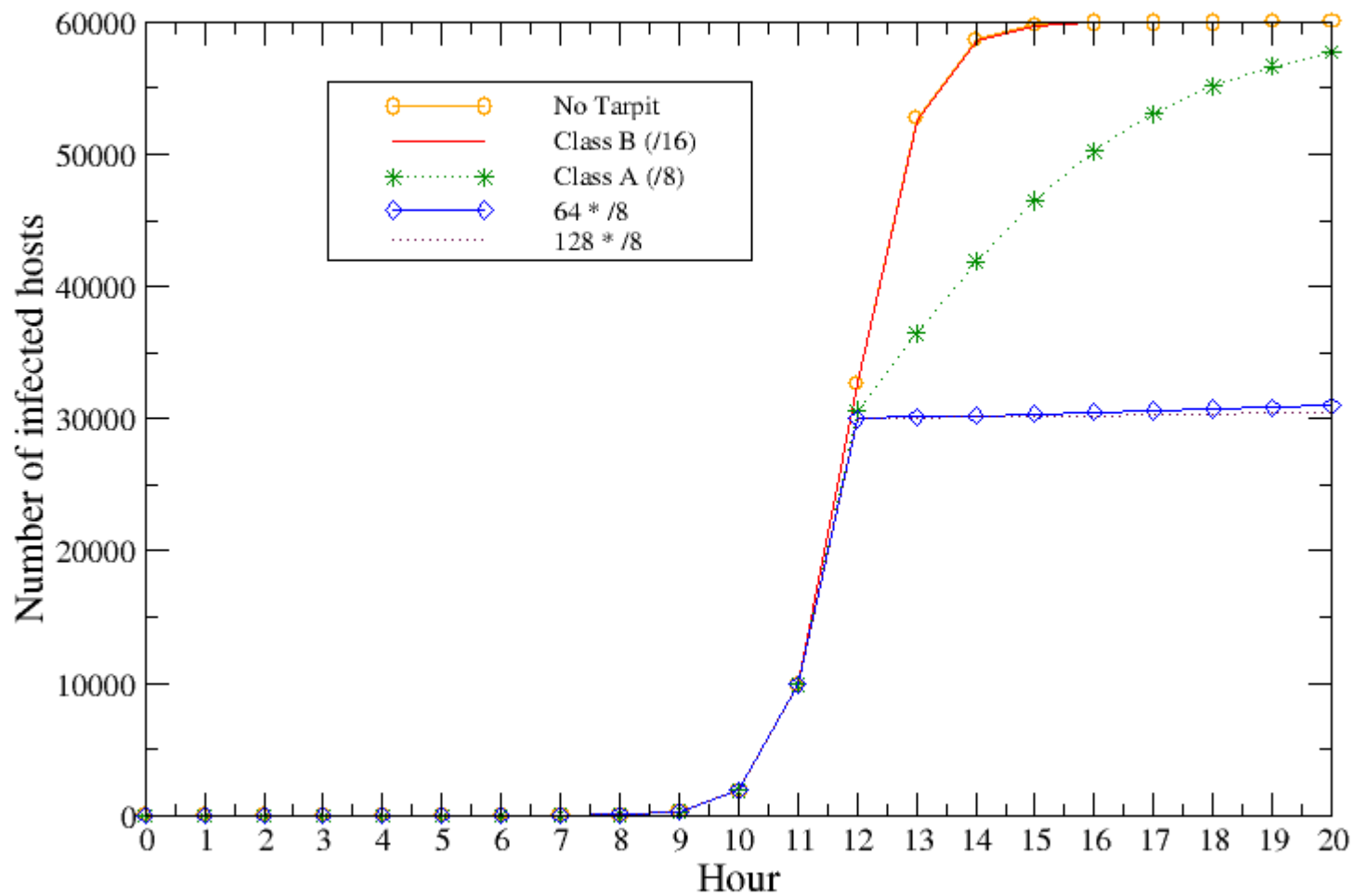
# iSink Active Responses

## Inbound Flows, Passive Control Groups

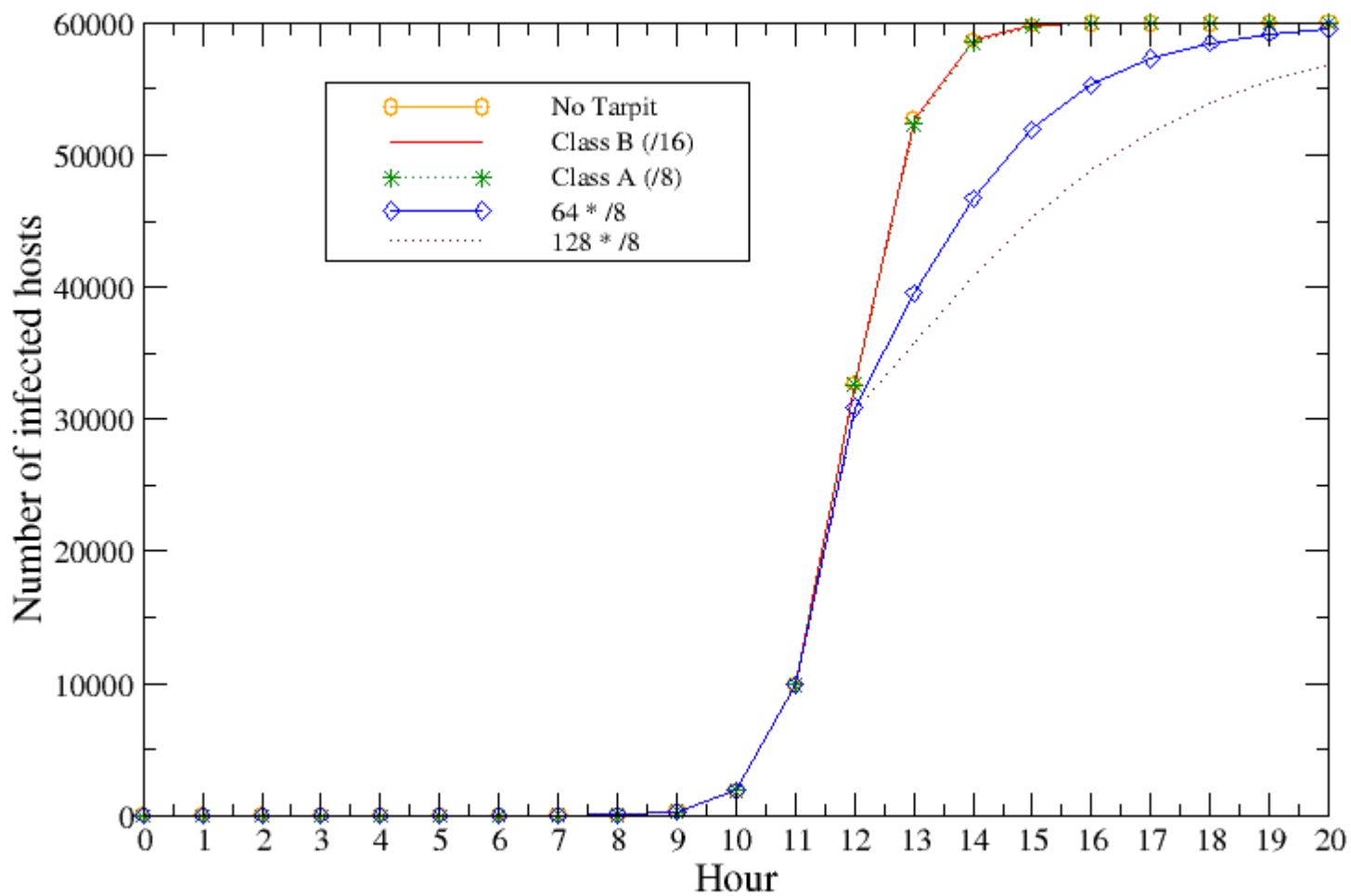
WAIL Class-A Network Sink Passive on 10 "/16" Subnets



# iSink Tarpit Effectiveness



# iSink Tarpit Effectiveness



# iSink Results

- Our iSinks are a novel application of the Click modular router and argus flow-based measurement tool.
- iSinks based on commodity PCs can handle ~20,000 connection attempts per second, well over the average amount of unsolicited traffic received by an "unused" class A network.
- Packet sampling can be used to effectively produce a list of the "top talker" sources.



# iSink Discussion

- „ How stealthy must an iSink be to get valid results?  
Is the iSink vulnerable to obfuscation?
- „ Legal Issues?
- „ What about IPv6?
  - much larger, much more sparsely populated space
- „ Enhance the global BGP to supported advertising short-cut routes to deliver unwanted traffic to iSinks? (eg. Barry Greene on advertising dark IP)
- „ Can Internet appliances and middleboxes act as a sink for unused transport endpoints? (eg. a tarpitting PAT router)

# iSink Future Work

- Expand monitoring to other networks
  - Are our results portable to other prefixes?
- Add NIDS-like signature recognition and alert capability
- Explore the efficacy of packet sampling
  - Can packet-sampling iSinks detect rare events?